# FURTHER INVESTIGATIONS ON HARRIS ALGORITHM 

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#### Abstract

In this note we will make further computational improvements of Harris algorithm [2, 12]. We improve speed using the technique of least absolute remainder [1]. Numerical experiment give us confidence that we receive new enhanced algorithm.


Key words: Euclidean algorithm, Harris algorithm, hybrid algorithm, least absolute remainder.

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## 1. Introduction

Harris algorithm is well known hybrid iteration process which compute Greatest common divisor of two natural numbers $a$ and $b$. In many classical and recent books and papers the Euclidean algorithm is well described, see [2]-[10] and [28]-[39]. Using symmetry properties of Euclidean iteration process we receive some computational benefits [11]-[27].

For testing purposes we will use the following computer: processor - Intel(R) Core(TM) i7-6700HQ CPU 2.60GHz, $2592 \mathrm{Mhz}, 4$ Core(s), 8 Logical Processor(s), RAM 16 GB, Microsoft Windows 10 Enterprise x64, Microsoft Visual C\# $2017 \times 64$.

## 2. Main Results

We present new iteration process, which improve Harris algorithm:

## Algorithm 1.

```
int \(\mathrm{g}=0\);
if \(((\mathrm{a} \& 1)==0 \& \&(\mathrm{~b} \& 1)==0)\)
do \(\{\mathrm{a} \gg=1 ; \mathrm{b} \gg=1 ; \mathrm{g}++;\}\)
while \(((\mathrm{a} \& 1)==0 \& \&(\mathrm{~b} \& 1)==0)\);
\(u=a ; v=b\)
while \(((\mathrm{u} \& 1)==0) \mathrm{u} \gg=1\);
```

while ( $(\mathrm{v} \& 1)==0) \mathrm{v} \gg=1$;
if $(u>v)$ do $\{u \%=v$;
if $(u<1)\{\operatorname{gcd}=\mathrm{v} \ll \mathrm{g}$; break; \}
if $((u \& 1)==0)$
$\{$ do $u \gg=1$; while $((u \& 1)==0)$;
if $(\mathrm{u}==1)\{\operatorname{gcd}=\mathrm{u} \ll \mathrm{g}$; break; $\}\}$
else $\{\mathrm{ar}=\mathrm{v}-\mathrm{u}$;
if ( $u>a r$ )
$\{\mathrm{u}=\mathrm{ar} ;$
do $u \gg=1$; while $((\mathrm{u} \& 1)==0)$;
if $(\mathrm{u}==1)\{\operatorname{gcd}=\mathrm{u} \ll \mathrm{g}$; break; $\}\}\}$
$\mathrm{v} \%=\mathrm{u}$;
if $(\mathrm{v}<1)\{$ gcd $=\mathrm{u} \ll \mathrm{g}$; break; \}
if $((\mathrm{v} \& 1)==0)$
$\{$ do $\mathrm{v} \gg=1$; while $((\mathrm{v} \& 1)==0)$;
if $(\mathrm{v}==1)\{$ gcd $=\mathrm{v} \ll \mathrm{g}$; break; $\}\}$
else $\{\mathrm{ar}=\mathrm{u}-\mathrm{v}$;
if ( $\mathrm{v}>\mathrm{ar}$ )
$\{\mathrm{v}=\mathrm{ar} ;$
do $\mathrm{v} \gg=1$; while $((\mathrm{v} \& 1)==0)$;
if $(\mathrm{v}==1)\{\operatorname{gcd}=\mathrm{v} \ll \mathrm{g}$; break; $\}\}\}$
\} while (true);
else do $\{\mathrm{v} \%=\mathrm{u}$;
if $(\mathrm{v}<1)$ \{ gcd $=\mathrm{u} \ll \mathrm{g}$; break; \}
if $((\mathrm{v} \& 1)==0)$
$\{$ do $\mathrm{v} \gg=1$; while $((\mathrm{v} \& 1)==0)$;
if $(\mathrm{v}==1)\{\mathrm{gcd}=\mathrm{v} \ll \mathrm{g}$; break; $\}\}$
else $\{\mathrm{ar}=\mathrm{u}-\mathrm{v}$;
if ( $\mathrm{v}>$ ar)
$\{\mathrm{v}=$ ar;
do $\mathrm{v} \gg=1$; while $((\mathrm{v} \& 1)==0)$;
if $(\mathrm{v}==1)\{\operatorname{gcd}=\mathrm{v} \ll \mathrm{g} ;$ break; $\}\}\}$
u $\%=\mathrm{v}$;
if $(\mathrm{u}<1)\{\mathrm{gcd}=\mathrm{v} \ll \mathrm{g}$; break; \}
if $((u \& 1)==0)$

```
\(\{\) do \(u \gg=1\); while \(((u \& 1)==0)\);
if \((\mathrm{u}==1)\{\mathrm{gcd}=\mathrm{u} \ll \mathrm{g}\); break; \(\}\}\)
else \(\{\mathrm{ar}=\mathrm{v}-\mathrm{u}\);
if ( \(u>a r\) )
\(\{u=\) ar;
do \(u \gg=1\); while \(((\mathrm{u} \& 1)==0)\);
if \((\mathrm{u}==1)\{\operatorname{gcd}=\mathrm{u} \ll \mathrm{g}\); break; \(\}\}\}\)
\} while (true);
as well as its recursive version
```


## Algorithm 2.

static long Euclid(long u, long v, int g)
\{
long ar;
if $(\mathrm{u}>\mathrm{v})\{\mathrm{u} \%=\mathrm{v}$;
if $(\mathrm{u}<1)\{$ return $\mathrm{v} \ll \mathrm{g}$; \}
if $((u \& 1)==0)$
return $\operatorname{Euclid}(\mathrm{u} \gg 1, \mathrm{v}, \mathrm{g})$;
else $\{$ if $(u==1)\{$ return $u \ll g$; \}
ar $=\mathrm{v}-\mathrm{u}$;
if ( $u>$ ar)
$\{\mathrm{u}=\mathrm{ar} ;$
if $((u \& 1)==0)$
return $\operatorname{Euclid}(u \gg 1, \mathrm{v}, \mathrm{g}) ;\}\}\}$
else $\{\mathrm{v} \%=\mathrm{u}$;
if $(\mathrm{v}<1)$ \{ return $\mathrm{u} \ll \mathrm{g}$; \}
if $((\mathrm{v} \& 1)==0)$
return $\operatorname{Euclid}(u, v \gg 1, g)$;
else $\{$ if $(\mathrm{v}==1)\{$ return $\mathrm{v} \ll \mathrm{g}$; \}
ar $=u-v$;
if ( $\mathrm{v}>$ ar)
$\{\mathrm{v}=\mathrm{ar}$;
if $((\mathrm{v} \& 1)==0)$
return $\operatorname{Euclid}(u, v \gg 1, g) ;\}\}\}$
return $\operatorname{Euclid}(u, v, g)$;
\}

The recursive function should be called by:
int $\mathrm{g}=0$;
if $((\mathrm{a} \& 1)==0 \& \&(\mathrm{~b} \& 1)==0)$
do $\{\mathrm{a} \gg=1 ; \mathrm{b} \gg=1 ; \mathrm{g}++;\}$
while $((\mathrm{a} \& 1)==0 \& \&(\mathrm{~b} \& 1)==0)$;
$\mathrm{u}=\mathrm{a} ; \mathrm{v}=\mathrm{b}$;
while $((\mathrm{u} \& 1)==0) \mathrm{u} \gg=1$;
while $((\mathrm{v} \& 1)==0) \mathrm{v} \gg=1$;
gcd $=\operatorname{Euclid}(\mathrm{u}, \mathrm{v}, \mathrm{g})$;

## Numerical Example.

For testing purposes of Algorithms 1 and 2 we will use the following main function:
long $\mathrm{a}, \mathrm{b}, \mathrm{gcd}, \mathrm{d} 1=0, \mathrm{u}, \mathrm{v}$;
for (int $\mathrm{i}=1 ; \mathrm{i}<100000001 ; \mathrm{i}++$ ) $\{\mathrm{a}=\mathrm{i} ; \mathrm{b}=200000002$ - i ;
//here are placed the source code of algorithm 1 and
//calling of recursive algorithm 2
$\mathrm{d} 1+=\mathrm{gcd} ;$
\}
Console.WriteLine(d1);
CPU time results are:
CPU time of Algorithm 1 is: $\mathbf{2 6 . 7 9 9}$ seconds.
CPU time of Algorithm 2 is: 42.989 seconds.
For the same numerical example Harris algorithms [2, 12] gave the following results - iterative 31.620 seconds and recursive 68.119 seconds.

## 3. Conclusion

We give how in Harris algorithm can be implemented the technique of least absolute remainder and this leads to computational speed improvements.

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